# Chapter 13 Optimization and Matching

## Section 13.1

### Definitions

weight wt(e), weighted graph, distance measure in graphs, shortest path, greedy algorithm

#### Algorithm

Dijkstra's Shortest-Path Algorithm

#### Exercises

1, 2, 5

## Section 13.2

## Definitions

optimal (minimal) spanning tree

### Algorithms

Kruskal's AlgorithmTh 13.1 $O(|V|^2 \log_2(|V|))$ Prim's AlgorithmTh 13.2 $O(|V|^2)$ 

#### Exercises

1, 2, 5, 6

## Section 11.3

### Definitions

network (transport network), source, sink, capacity, flow, saturated, cut, value of the flow val(f), capacity of cut  $c(P, \bar{P})$ 

### Theorem

 $\begin{array}{ll} 13.3 & C = (P,\bar{P}) \text{ any cut. Then } c(P,\bar{P}) \geq val(f).\\ \text{C13.1} & \text{value of flow from the source} = \text{value of flow into the sink.}\\ 13.4 & \min c(P,\bar{P}) = val(f) \end{array}$ 

### Exercises

1, 2, 4, 6, 7

### Section 11.4

### Definitions

assignment problem, matching, complete matching, maximal matching, perfect matching, deficiency of, deficiency of graph G,  $\delta(G)$ , permutation matrix, doubly stochastic

### Theorem

G bipartite, partitioned as  $X\cup Y.$ 

- 13.5
- C13.3
- Complete matching of X into Y  $\Leftrightarrow \forall A \subseteq X : |A| \leq |\{y \in Y | \exists x \in A : x \text{ adjacent to } y\}|$  $\exists k \in \mathbb{Z}^+ \forall x \in X, y \in Y : deg(x) \geq k \geq deg(y) \Rightarrow \text{Complete matching of X into Y}.$ The maximum number of vertices in X that can be matched with those in Y is  $|X| \delta(G)$ . 13.6

#### Exercises

3, 5, 7, 13, 14Supplementary 2, 7, 8 (Birkhoff-von Neumann Theorem)